### Introductory Programming Recursion

#### IMM, DTU

a. Parts of this material are inspired by/originate from a course at ITU developed by Niels Hallenberg and Peter Sestoft on the basis of a course at KVL developed by Morten Larsen and Peter Sestoft. Translated into English and edited by Anne Haxthausen.

© Haxthausen and Sestoft, IMM/DTU, 19th November 2002

02100+02115+02199+02312 Introductory Programming

Page 11-1

## Recursion: methods that call themselves

A method can invoke another method

A method can also invoke itself. Such a method is said to be recursive.

© Haxthausen and Sestoft, IMM/DTU, 19th November 2002

02100+02115+02199+02312 Introductory Programming

mming Page 11-2

### **Example: factorial function**

The function n!, pronounced n factorial, gives the number of permutations of n items. E.g. 3 persons can sit on a bench in  $3!=3\cdot 2\cdot 1=6$  different permutations.

The factorial function n! can be defined as follows

$$n! = n \cdot (n-1) \cdot \ldots \cdot 1$$

However, the factorial function can also be defined recursively:

$$= \begin{cases} 1 & \text{if } n = 0 \\ n \cdot (n-1)! & \text{if } n > 0 \end{cases}$$

The first definition is most easily implemented as a loop (i.e. as an iterative method).

The second definition is most easily implemented as a method that calls itself (i.e. as a recursive method).

©Haxthausen and Sestoft, IMM/DTU, 19th November 2002

02100+02115+02199+02312 Introductory Programming

Page 11-3

## Iterative and recursive calculations of n!

```
public class Factorial {
   public static void main(String args[]) {
    int n = Integer.parseInt(args[0]);
    System.out.println(n + "! is " + ifac(n));
    System.out.println(n + "! is " + rfac(n));
}

static int ifac(int n) { // Iterative factorial function int result = 1;
   for (int i=n; i >= 1; i=i-1)
        result = result * i;
        return result;
}

static int rfac(int n) { // Recursive factorial function if (n == 0) return 1;
        else
        return n * rfac(n - 1);
}
```

©Haxthausen and Sestoft, IMM/DTU, 19th November 2002

U, 19th November 2002 02100+02115+0219

02100+02115+02199+02312 Introductory Programming

Page 11-4

## How rfac(3) is evaluated in the machine

```
ω
             ω
                          3 * (2 * (1 * rfac(0)))
                                                      3 * rfac(2)
                                                                   rfac(3)
                                        * (2 * rfac(1))
```

completed expressions During the evaluation, the recursive method invocations require allocation of memory for not

Recursive method invocations are a little slower than repetitions in a loop.

©Haxthausen and Sestoft, IMM/DTU, 19th November 2002

02100+02115+02199+02312 Introductory Programming

Page 11-5

©Haxthausen and Sestoft, IMM/DTU, 19th November 2002

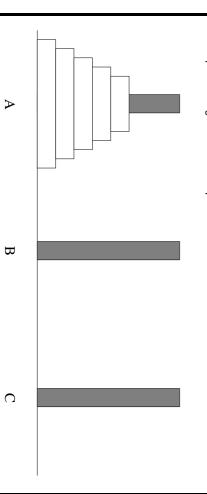
# Some problems are more elegantly solved using recursive methods

Example: the towers of Hanoi

All disks should be moved from tower A to tower B.

Only move one disk at a time

Never place a larger disk on the top of a smaller disk



© Haxthausen and Sestoft, IMM/DTU, 19th November 2002

02100+02115+02199+02312 Introductory Programming

Page 11-6

## Example: the towers of Hanoi, algorithm

Move  $\boldsymbol{n}$  disks from tower A to tower B via an extra tower C:

- 1. Move the n-1 topmost disks from tower A to tower C.
- 2. Move a disk from tower A to tower B. 3. Move the n-1 topmost disks from tower C to tower B.

To make step 1 use the recipe replacing n with n-1, 'B' with 'C' and 'C' with 'B':

Move n-1 disks from tower A to tower C via an extra tower B:

- 1. Move the n-2 topmost disks from tower A to tower B. 2. Move a disk from tower A to tower C. 3. Move the n-2 topmost disks from tower B to tower C.

Similarly, to make step 3 replace 'n' with 'n-1', 'A' with 'C' and 'C' with 'A':

Move n-1 disks from tower C to tower B via an extra tower A: 1. Move the n-2 topmost disks from tower C to tower A.

- 2. Move a disk from tower C to tower B. 3. Move the n-2 topmost disks from tower A to tower B.

In order to move zero disks you do not need to do anything

02100+02115+02199+02312 Introductory Programming

Page 11-7

### Hanoi – example with 3 disks

- 1. Move disk 1, 2 from A to C via B
- (a) Move disk 1 from A to B
- (b) Move disk 2 from A to C
- (c) Move disk 1 from B to C
- 2. Move disk 3 from A to B
- 3. Move disk 1, 2 from C to B via A
- (a) Move disk 1 from C to A
- (c) Move disk 1 from A to B (b) Move disk 2 from C to B

© Haxthausen and Sestoft, IMM/DTU, 19th November 2002

02100+02115+02199+02312 Introductory Programming

## Example: 'the towers of Hanoi' in Java (Hanoi. java)

© Haxthausen and Sestoft, IMM/DTU, 19th November 2002

02100+02115+02199+02312 Introductory Programming

Page 11-9

### Main idea in recursion

Some problems can be solved by solving one or more subproblems of the same kind

E.g. the towers of Hanoi: move a stack of  $\boldsymbol{n}$  disks by

- ullet moving a stack of n-1 disks,
- moving 1 disk, and
- ullet moving a stack of n-1 disks

Such a strategy for problem solving is called 'divide-and-conquer' (Danish: 'Del og hersk').

It is well-suited for implementation with recursive methods.

The strategy only works if:

- each subproblem is simpler than the original one, and
- it terminates with a trivial subproblem (e.g. move a stack of 0 disks)

© Haxthausen and Sestoft, IMM/DTU, 19th November 2002

02100+02115+02199+02312 Introductory Programming

Page 11-10